## Structured Flowcharts

```
grammar StructuredFlowDiagrams
  nonterminal D(2), Stmt(2), S(0);
   terminal begin(1), stop(1), action(2), pred(3);
   start S:
  S()
             ::= begin(x) D(x,y) stop(y)
                                             [ init
  D(x,y)
             ::= Stmt(x,y)
                                             [ single ]
              | Stmt(x,z) D(z,y)
                                             [ seq
  Stmt(x,y) ::= action(x,y)
                                             [action]
              | pred(x,u,v) D(u,y) D(v,y)
                                             [alt
                                                      1
              | pred(x,z,y) D(z,x)
                                             [ while
                                                    ]
end
```

Structured flowcharts are not PSR because the CFA has conflicts (see below). This CFA is somewhat special because parameters, different from all other examples considered so far, can not only refer to a single node of the host graph, but to an arbitrary number. This is the case for formal parameter  $\mathbf{a}$  in states  $q_2(\mathbf{a}, \mathbf{b}^*, \mathbf{c}, \mathbf{d})$  and  $q_{10}(\mathbf{a}, \mathbf{b}, \mathbf{c}^*)$ . To indicate these multi parameters, they are marked with an asterisk. The multi value is obtained by two transitions leaving from  $q_2(\mathbf{a}, \mathbf{b}^*, \mathbf{c}, \mathbf{d})$  where the value of the target parameter is determined by expressions  $\mathbf{a} \cdot \mathbf{b}$  and  $\mathbf{b} \cdot \mathbf{a}$ , respectively. These expressions mean that the value of  $\mathbf{b}$ , i.e., a list of nodes, is concatenated with the value of  $\mathbf{a}$ , i.e., a single node. That way, a longer list of host nodes is obtained and assigned to to corresponding formal parameter. Whenever such a multi parameter is used, any of its assigned host nodes can be used, for instance in the second transition leaving  $q_2(\mathbf{a}, \mathbf{b}^*, \mathbf{c}, \mathbf{d})$ , indicated by the condition  $\exists i: n_0 = \mathbf{c}, n_1 = \mathbf{b}_i$ .  $\mathbf{b}_i$  stands for the node at position i of the list being the value of  $\mathbf{b}$ . But as this grammar is not PSR, this has no effect on a parser.

## Characteristic finite automaton(CFA):

**Note:** State  $q_6$  has conflicts

## State $q_0(a, b)$

$$\mathsf{S}() o \mathtt{f .} \mathsf{begin}(m{a}) \, \mathsf{D}(m{a}, m{b}) \, \mathsf{stop}(m{b})$$

$$\frac{\mathsf{begin}(n_0)}{n_0 = \boldsymbol{a}} \to q_1(n_0, \boldsymbol{b})$$

#### State $q_1(a,b)$

$$\begin{array}{|c|c|}\hline S() & \rightarrow \mathsf{begin}(\boldsymbol{a}) \cdot \mathsf{D}(\boldsymbol{a},\boldsymbol{b}) \, \mathsf{stop}(\boldsymbol{b}) \\ \hline \mathsf{D}(\boldsymbol{a},\boldsymbol{b}) & \rightarrow \cdot \mathsf{Stmt}(\boldsymbol{a},\boldsymbol{b}) \\ \mathsf{D}(\boldsymbol{a},\boldsymbol{b}) & \rightarrow \cdot \mathsf{Stmt}(\boldsymbol{a},n_1) \, \mathsf{D}(n_1,\boldsymbol{b}) \\ \mathsf{Stmt}(\boldsymbol{a},\boldsymbol{b}) & \rightarrow \cdot \mathsf{action}(\boldsymbol{a},\boldsymbol{b}) \\ \mathsf{Stmt}(\boldsymbol{a},\boldsymbol{b}) & \rightarrow \cdot \mathsf{pred}(\boldsymbol{a},n_2,n_3) \, \mathsf{D}(n_2,\boldsymbol{b}) \, \mathsf{D}(n_3,\boldsymbol{b}) \\ \mathsf{Stmt}(\boldsymbol{a},\boldsymbol{b}) & \rightarrow \cdot \mathsf{pred}(\boldsymbol{a},n_4,\boldsymbol{b}) \, \mathsf{D}(n_4,\boldsymbol{a}) \\ \mathsf{Stmt}(\boldsymbol{a},n_5) & \rightarrow \cdot \mathsf{action}(\boldsymbol{a},n_5) \\ \mathsf{Stmt}(\boldsymbol{a},n_6) & \rightarrow \cdot \mathsf{pred}(\boldsymbol{a},n_7,n_8) \, \mathsf{D}(n_7,n_6) \, \mathsf{D}(n_8,n_6) \\ \mathsf{Stmt}(\boldsymbol{a},n_9) & \rightarrow \cdot \mathsf{pred}(\boldsymbol{a},n_{10},n_9) \, \mathsf{D}(n_{10},\boldsymbol{a}) \\ \end{array}$$

$$\begin{array}{c} \begin{array}{c} \mathbf{D}(n_0,n_1) \\ \hline n_0 = \boldsymbol{a}, n_1 = \boldsymbol{b} \\ \hline \mathbf{Stmt}(n_0,n_1) \\ \hline n_0 = \boldsymbol{a}, n_1 \uparrow \\ \hline \mathbf{Stmt}(n_0,n_1) \\ \hline n_0 = \boldsymbol{a}, n_1 \uparrow \\ \hline \\ \mathbf{Stmt}(n_0,n_1) \\ \hline n_0 = \boldsymbol{a}, n_1 = \boldsymbol{b} \\ \hline \mathbf{action}(n_0,n_1) \\ \hline n_0 = \boldsymbol{a}, n_1 \uparrow \\ \hline \mathbf{action}(n_0,n_1) \\ \hline n_0 = \boldsymbol{a}, n_1 \uparrow \\ \hline \mathbf{action}(n_0,n_1) \\ \hline n_0 = \boldsymbol{a}, n_1 = \boldsymbol{b} \\ \hline \\ \mathbf{pred}(n_0,n_1,n_2) \\ \hline n_0 = \boldsymbol{a}, n_1 \uparrow, n_2 \uparrow \\ \hline \\ \mathbf{pred}(n_0,n_1,n_2) \\ \hline \\ \mathbf{n}_0 = \boldsymbol{a}, n_1 \uparrow, n_2 \uparrow \\ \hline \\ \mathbf{pred}(n_0,n_1,n_2) \\ \hline \\ \mathbf{n}_0 = \boldsymbol{a}, n_1 \uparrow, n_2 = \boldsymbol{b} \\ \hline \end{array} \right. \qquad q_2(n_0,n_1)$$

#### State $q_2(a,b)$

ig| Stmt $(oldsymbol{a},oldsymbol{b}) o$  action $(oldsymbol{a},oldsymbol{b})$  . [action]

### State $q_3(a, b^*, c, d)$

```
\mathsf{Stmt}(m{a},m{b}) 	o \mathsf{pred}(m{a},m{c},m{d}) . \mathsf{D}(m{c},m{b})\,\mathsf{D}(m{d},m{b})
\mathsf{Stmt}(m{a},m{d}) 	o \mathsf{pred}(m{a},m{c},m{d}) . \mathsf{D}(m{c},m{a})
\mathsf{Stmt}(\boldsymbol{a}, n_1) \to \mathsf{pred}(\boldsymbol{a}, \boldsymbol{c}, \boldsymbol{d}) \cdot \mathsf{D}(\boldsymbol{c}, n_1) \mathsf{D}(\boldsymbol{d}, n_1)
\mathsf{D}(\boldsymbol{c}, \boldsymbol{a})

ightarrow • Stmt(oldsymbol{c},oldsymbol{a})
                                      \rightarrow • Stmt(\boldsymbol{c}, n_2) D(n_2, \boldsymbol{a})
\mathsf{D}(\boldsymbol{c},\boldsymbol{a})
\mathsf{D}(\boldsymbol{c},n_3)
                                      \rightarrow \bot \mathsf{Stmt}(\boldsymbol{c}, n_3)
\mathsf{D}(\boldsymbol{c},n_4)

ightarrow • Stmt(oldsymbol{c},n_5) D(n_5,n_4)
\mathsf{D}(\boldsymbol{c},\boldsymbol{b})

ightarrow ullet Stmt(oldsymbol{c},oldsymbol{b})
\mathsf{D}(\boldsymbol{c},\boldsymbol{b})
                                      \rightarrow Stmt(\boldsymbol{c}, n_6) D(n_6, \boldsymbol{b})
\mathsf{Stmt}(oldsymbol{c},oldsymbol{a}) 	o \mathtt{action}(oldsymbol{c},oldsymbol{a})
\mathsf{Stmt}(\boldsymbol{c}, \boldsymbol{a}) \quad 	o \, \mathsf{pred}(\boldsymbol{c}, n_7, n_8) \, \mathsf{D}(n_7, \boldsymbol{a}) \, \mathsf{D}(n_8, \boldsymbol{a})
\mathsf{Stmt}(oldsymbol{c},oldsymbol{a}) 	o \mathtt{.pred}(oldsymbol{c},n_9,oldsymbol{a})\,\mathsf{D}(n_9,oldsymbol{c})
\mathsf{Stmt}(oldsymbol{c},n_{10}) 	o \mathtt{Laction}(oldsymbol{c},n_{10})
\mathsf{Stmt}(\boldsymbol{c}, n_{11}) \to \mathsf{pred}(\boldsymbol{c}, n_{12}, n_{13}) \, \mathsf{D}(n_{12}, n_{11}) \, \mathsf{D}(n_{13}, n_{11})
\mathsf{Stmt}(\boldsymbol{c}, n_{14}) \to \mathsf{pred}(\boldsymbol{c}, n_{15}, n_{14}) \, \mathsf{D}(n_{15}, \boldsymbol{c})
\mathsf{Stmt}(\boldsymbol{c}, \boldsymbol{b})

ightarrow a action(oldsymbol{c},oldsymbol{b})
\mathsf{Stmt}(\boldsymbol{c}, \boldsymbol{b})
                                    \rightarrow \mathsf{pred}(\boldsymbol{c}, n_{16}, n_{17}) \, \mathsf{D}(n_{16}, \boldsymbol{b}) \, \mathsf{D}(n_{17}, \boldsymbol{b})
\mathsf{Stmt}(\boldsymbol{c}, \boldsymbol{b})

ightarrow ullet pred(oldsymbol{c}, n_{18}, oldsymbol{b}) \, \mathsf{D}(n_{18}, oldsymbol{c})
```

## State $q_4(a, b, c)$

$Stmt(m{a},m{b})$	$ ightarrow$ pred $(oldsymbol{a},oldsymbol{c},oldsymbol{b})$ . $D(oldsymbol{c},oldsymbol{a})$
$D(oldsymbol{c},oldsymbol{a})$	$ ightarrow$ . $Stmt(oldsymbol{c},oldsymbol{a})$
$D(oldsymbol{c},oldsymbol{a})$	$ ightarrow$ • Stmt $(oldsymbol{c},n_1)$ D $(n_1,oldsymbol{a})$
$Stmt(oldsymbol{c},oldsymbol{a})$	$ ightarrow$ , $action(oldsymbol{c},oldsymbol{a})$
$Stmt(oldsymbol{c},oldsymbol{a})$	$ ightarrow$ $ullet$ pred $(oldsymbol{c},n_2,n_3)D(n_2,oldsymbol{a})D(n_3,oldsymbol{a})$
$Stmt(oldsymbol{c},oldsymbol{a})$	$ ightarrow$ $ullet$ pred $(oldsymbol{c},n_4,oldsymbol{a})$ D $(n_4,oldsymbol{c})$
$Stmt(\boldsymbol{c}, n_5)$	$ ightarrow$ $ullet$ action $(oldsymbol{c},n_5)$
$Stmt(\boldsymbol{c},n_6)$	$ ightarrow$ • pred $(oldsymbol{c},n_7,n_8)D(n_7,n_6)D(n_8,n_6)$
$Stmt(oldsymbol{c},n_9)$	$ ightarrow$ $ullet$ pred $(oldsymbol{c},n_{10},n_{9})$ D $(n_{10},oldsymbol{c})$

$$\begin{array}{c} \begin{array}{c} D(n_0,n_1) \\ \hline n_0 = \boldsymbol{c}, n_1 = \boldsymbol{a} \\ \hline Stmt(n_0,n_1) \\ \hline n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline Stmt(n_0,n_1) \\ \hline \hline n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline Stmt(n_0,n_1) \\ \hline \hline n_0 = \boldsymbol{c}, n_1 = \boldsymbol{a} \\ \hline action(n_0,n_1) \\ \hline n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline action(n_0,n_1) \\ \hline n_0 = \boldsymbol{c}, n_1 = \boldsymbol{a} \\ \hline \hline n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline action(n_0,n_1) \\ \hline n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline oped(n_0,n_1,n_2) \\ \hline \hline n_0 = \boldsymbol{c}, n_1 \uparrow, n_2 \uparrow \\ \hline pred(n_0,n_1,n_2) \\ \hline \hline n_0 = \boldsymbol{c}, n_1 \uparrow, n_2 \uparrow \\ \hline oped(n_0,n_1,n_2) \\ \hline \hline oped(n_0,n_1,n_2) \\ \hline oped(n_0,n_1,n_2$$

# State $q_5(a, b)$

 $\mathsf{D}(oldsymbol{a},oldsymbol{b}) o \mathsf{Stmt}(oldsymbol{a},oldsymbol{b})$  . [single]

## State $q_6(a, b, c^*)$

$$\begin{array}{c|cccc} D(\boldsymbol{a},\boldsymbol{b}) & \rightarrow \operatorname{Stmt}(\boldsymbol{a},\boldsymbol{b}) \cdot & [single] \\ D(\boldsymbol{a},n_1) & \rightarrow \operatorname{Stmt}(\boldsymbol{a},\boldsymbol{b}) \cdot \operatorname{D}(\boldsymbol{b},n_1) \\ D(\boldsymbol{a},\boldsymbol{c}) & \rightarrow \operatorname{Stmt}(\boldsymbol{a},\boldsymbol{b}) \cdot \operatorname{D}(\boldsymbol{b},\boldsymbol{c}) \\ \hline D(\boldsymbol{b},n_2) & \rightarrow \cdot \operatorname{Stmt}(\boldsymbol{b},n_2) \\ D(\boldsymbol{b},n_3) & \rightarrow \cdot \operatorname{Stmt}(\boldsymbol{b},n_4) \operatorname{D}(n_4,n_3) \\ D(\boldsymbol{b},\boldsymbol{c}) & \rightarrow \cdot \operatorname{Stmt}(\boldsymbol{b},\boldsymbol{c}) \\ D(\boldsymbol{b},\boldsymbol{c}) & \rightarrow \cdot \operatorname{Stmt}(\boldsymbol{b},n_5) \operatorname{D}(n_5,\boldsymbol{c}) \\ \operatorname{Stmt}(\boldsymbol{b},n_6) & \rightarrow \cdot \operatorname{action}(\boldsymbol{b},n_6) \\ \operatorname{Stmt}(\boldsymbol{b},n_7) & \rightarrow \cdot \operatorname{pred}(\boldsymbol{b},n_8,n_9) \operatorname{D}(n_8,n_7) \operatorname{D}(n_9,n_7) \\ \operatorname{Stmt}(\boldsymbol{b},n_{10}) & \rightarrow \cdot \operatorname{pred}(\boldsymbol{b},n_{11},n_{10}) \operatorname{D}(n_{11},\boldsymbol{b}) \\ \operatorname{Stmt}(\boldsymbol{b},\boldsymbol{c}) & \rightarrow \cdot \operatorname{action}(\boldsymbol{b},\boldsymbol{c}) \\ \operatorname{Stmt}(\boldsymbol{b},\boldsymbol{c}) & \rightarrow \cdot \operatorname{pred}(\boldsymbol{b},n_{12},n_{13}) \operatorname{D}(n_{12},\boldsymbol{c}) \operatorname{D}(n_{13},\boldsymbol{c}) \\ \operatorname{Stmt}(\boldsymbol{b},\boldsymbol{c}) & \rightarrow \cdot \operatorname{pred}(\boldsymbol{b},n_{14},\boldsymbol{c}) \operatorname{D}(n_{14},\boldsymbol{b}) \end{array}$$

$$\begin{array}{c} \begin{array}{c} D(n_0,n_1) \\ \hline n_0 = \boldsymbol{b}, n_1 \uparrow \\ \hline D(n_0,n_1) \\ \hline \exists i: n_0 = \boldsymbol{b}, n_1 = \boldsymbol{c}_i \end{array} & \rightarrow q_8(\boldsymbol{a},n_1,n_0) \\ \hline Stmt(n_0,n_1) \\ \hline n_0 = \boldsymbol{b}, n_1 \uparrow \\ \hline Stmt(n_0,n_1) \\ \hline \exists i: n_0 = \boldsymbol{b}, n_1 = \boldsymbol{c}_i \end{array} & \rightarrow q_6(n_0,n_1,\boldsymbol{c}) \\ \hline Stmt(n_0,n_1) \\ \hline \exists i: n_0 = \boldsymbol{b}, n_1 = \boldsymbol{c}_i \\ \hline \text{action}(n_0,n_1) \\ \hline n_0 = \boldsymbol{b}, n_1 \uparrow \\ \hline \text{action}(n_0,n_1) \\ \hline \exists i: n_0 = \boldsymbol{b}, n_1 = \boldsymbol{c}_i \end{array} & \rightarrow q_2(n_0,n_1) \\ \hline \text{pred}(n_0,n_1,n_2) \\ \hline n_0 = \boldsymbol{b}, n_1 \uparrow, n_2 \uparrow \\ \hline \text{pred}(n_0,n_1,n_2) \\ \hline \hline n_0 = \boldsymbol{b}, n_1 \uparrow, n_2 \uparrow \\ \hline \text{pred}(n_0,n_1,n_2) \\ \hline \hline \exists i: n_0 = \boldsymbol{b}, n_1 \uparrow, n_2 = \boldsymbol{c}_i \end{array} & q_4(n_0,n_2,n_1) \end{array}$$

#### Conflicts:

• shift  $pred(b, \circ, \circ)$ , reduce single, shift  $action(b, \circ)$ 

#### State $q_7(a, b, c)$

2 tate 47 (a, 5, 5)		
$D(oldsymbol{a},oldsymbol{b})$	$ ightarrow Stmt(oldsymbol{a},oldsymbol{c})$ . $D(oldsymbol{c},oldsymbol{b})$	
$D(oldsymbol{c},oldsymbol{b})$	$ ightarrow$ . $Stmt(oldsymbol{c},oldsymbol{b})$	
$D(oldsymbol{c},oldsymbol{b})$	$ ightarrow$ $ullet$ Stmt $(oldsymbol{c},n_1)$ D $(n_1,oldsymbol{b})$	
$Stmt(oldsymbol{c},n_2)  o \mathtt{ullet} action(oldsymbol{c},n_2)$		
$Stmt(\boldsymbol{c},n_3)$	$0  ightarrow ullet$ pred $(oldsymbol{c}, n_4, n_5)  D(n_4, n_3)  D(n_5, n_3)$	
$Stmt(\boldsymbol{c},n_6)$	$0  ightarrow ullet$ pred $(oldsymbol{c}, n_7, n_6)  D(n_7, oldsymbol{c})$	
$Stmt(m{c}, m{b})$	$ ightarrow$ $ullet$ action $(oldsymbol{c},oldsymbol{b})$	
$Stmt(m{c}, m{b})$	$ ightarrow$ , $pred(oldsymbol{c}, n_8, n_9)D(n_8, oldsymbol{b})D(n_9, oldsymbol{b})$	
$Stmt(m{c}, m{b})$	$ ightarrow$ , $pred(oldsymbol{c}, n_{10}, oldsymbol{b})  D(n_{10}, oldsymbol{c})$	

$$\begin{array}{c} \begin{array}{c} D(n_0,n_1) \\ \hline n_0 = \boldsymbol{c}, n_1 = \boldsymbol{b} \\ \hline Stmt(n_0,n_1) \\ \hline n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline \\ N_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline \\ Stmt(n_0,n_1) \\ \hline \\ n_0 = \boldsymbol{c}, n_1 = \boldsymbol{b} \\ \hline \\ action(n_0,n_1) \\ \hline \\ n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline \\ action(n_0,n_1) \\ \hline \\ n_0 = \boldsymbol{c}, n_1 \uparrow \\ \hline \\ action(n_0,n_1) \\ \hline \\ n_0 = \boldsymbol{c}, n_1 = \boldsymbol{b} \\ \hline \\ pred(n_0,n_1,n_2) \\ \hline \\ n_0 = \boldsymbol{c}, n_1 \uparrow, n_2 \uparrow \\ \hline \\ pred(n_0,n_1,n_2) \\ \hline \\ n_0 = \boldsymbol{c}, n_1 \uparrow, n_2 \uparrow \\ \hline \\ pred(n_0,n_1,n_2) \\ \hline \\ n_0 = \boldsymbol{c}, n_1 \uparrow, n_2 = \boldsymbol{b} \\ \hline \end{array} \right. \\ q_8(\boldsymbol{a}, n_1, n_0)$$

### State $q_8(a, b, c)$

$$\mathsf{D}(oldsymbol{a},oldsymbol{b}) o \mathsf{Stmt}(oldsymbol{a},oldsymbol{c})\,\mathsf{D}(oldsymbol{c},oldsymbol{b})$$
 .  $[seq]$ 

## State $q_9(a, b)$

$$\mathsf{S}() o \mathsf{begin}(oldsymbol{a}) \, \mathsf{D}(oldsymbol{a}, oldsymbol{b})$$
 ,  $\mathsf{stop}(oldsymbol{b})$ 

$$\frac{\operatorname{\mathsf{stop}}(n_0)}{n_0 = \boldsymbol{b}} \to q_{12}(\boldsymbol{a}, n_0)$$

## State $q_{10}(a, b, c, d)$

$$\begin{array}{|c|c|}\hline \mathsf{Stmt}(\pmb{a},\pmb{b}) &\to \mathsf{pred}(\pmb{a},\pmb{c},\pmb{d})\,\mathsf{D}(\pmb{c},\pmb{b})\,\, \cdot\, \mathsf{D}(\pmb{d},\pmb{b})\\ \hline \mathsf{D}(\pmb{d},\pmb{b}) &\to \cdot\, \mathsf{Stmt}(\pmb{d},\pmb{b})\\ \mathsf{D}(\pmb{d},\pmb{b}) &\to \cdot\, \mathsf{Stmt}(\pmb{d},n_1)\,\mathsf{D}(n_1,\pmb{b})\\ \mathsf{Stmt}(\pmb{d},\pmb{b}) &\to \cdot\, \mathsf{action}(\pmb{d},\pmb{b})\\ \mathsf{Stmt}(\pmb{d},\pmb{b}) &\to \cdot\, \mathsf{pred}(\pmb{d},n_2,n_3)\,\mathsf{D}(n_2,\pmb{b})\,\mathsf{D}(n_3,\pmb{b})\\ \mathsf{Stmt}(\pmb{d},\pmb{b}) &\to \cdot\, \mathsf{pred}(\pmb{d},n_4,\pmb{b})\,\mathsf{D}(n_4,\pmb{d})\\ \mathsf{Stmt}(\pmb{d},\pmb{b}) &\to \cdot\, \mathsf{pred}(\pmb{d},n_4,\pmb{b})\,\mathsf{D}(n_4,\pmb{d})\\ \mathsf{Stmt}(\pmb{d},n_5) &\to \cdot\, \mathsf{action}(\pmb{d},n_5)\\ \mathsf{Stmt}(\pmb{d},n_6) &\to \cdot\, \mathsf{pred}(\pmb{d},n_7,n_8)\,\mathsf{D}(n_7,n_6)\,\mathsf{D}(n_8,n_6)\\ \mathsf{Stmt}(\pmb{d},n_9) &\to \cdot\, \mathsf{pred}(\pmb{d},n_{10},n_9)\,\mathsf{D}(n_{10},\pmb{d}) \end{array}$$

$$\begin{array}{l} \frac{\mathsf{D}(n_0,n_1)}{n_0 = \boldsymbol{d}, n_1 = \boldsymbol{b}} \to q_{13}(\boldsymbol{a},n_1,\boldsymbol{c},n_0) \\ \frac{\mathsf{Stmt}(n_0,n_1)}{n_0 = \boldsymbol{d},n_1 \uparrow} \to q_7(n_0,\boldsymbol{b},n_1) \\ \frac{\mathsf{Stmt}(n_0,n_1)}{n_0 = \boldsymbol{d},n_1 = \boldsymbol{b}} \to q_5(n_0,n_1) \\ \frac{\mathsf{action}(n_0,n_1)}{n_0 = \boldsymbol{d},n_1 \uparrow} \to q_2(n_0,n_1) \\ \frac{\mathsf{action}(n_0,n_1)}{n_0 = \boldsymbol{d},n_1 = \boldsymbol{b}} \to q_2(n_0,n_1) \\ \frac{\mathsf{pred}(n_0,n_1,n_2)}{n_0 = \boldsymbol{d},n_1 \uparrow,n_2 \uparrow} \to q_3(n_0,\boldsymbol{b},n_1,n_2) \\ \frac{\mathsf{pred}(n_0,n_1,n_2)}{n_0 = \boldsymbol{d},n_1 \uparrow,n_2 = \boldsymbol{b}} \to q_4(n_0,n_2,n_1) \end{array}$$

#### State $q_{11}(a,b,c)$

$$\mathsf{Stmt}(m{a},m{b}) o \mathsf{pred}(m{a},m{c},m{b})\,\mathsf{D}(m{c},m{a})$$
 .  $[\mathit{while}]$ 

### State $q_{12}(a,b)$

$$\mathsf{S}() o \mathsf{begin}(oldsymbol{a}) \, \mathsf{D}(oldsymbol{a}, oldsymbol{b}) \, \mathsf{stop}(oldsymbol{b}) \, \mathsf{l} \quad [\mathit{init}]$$

## State $q_{13}(a,b,c,d)$

$$\mathsf{Stmt}(m{a},m{b}) o \mathsf{pred}(m{a},m{c},m{d})\,\mathsf{D}(m{c},m{b})\,\mathsf{D}(m{d},m{b})$$
 .  $[alt]$