

# Deploying QKD in Network Layer VPN Infrastructures

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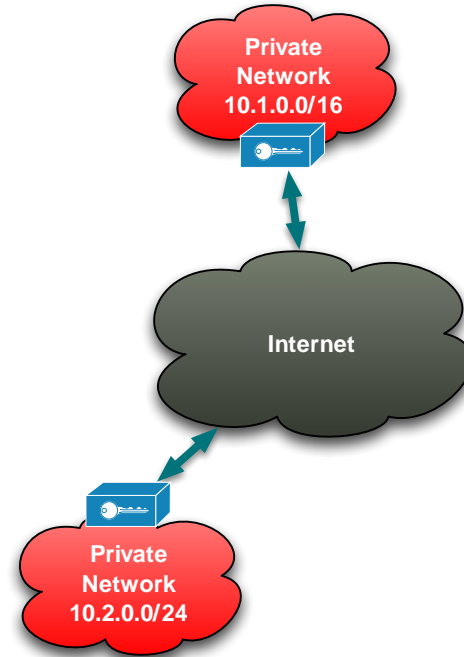
# Roadmap

1. **Network layer VPNs: Scenarios, requirements & current solutions**
2. **Challenge: How to defeat quantum computing attackers?**
3. **Emerging QKD standards: Overview and reflection**
4. **Integration of QKD & VPN Technology: An IT security guided approach**
5. **Conclusions**

# Network Layer VPN Infrastructures (1)

## Scenario:

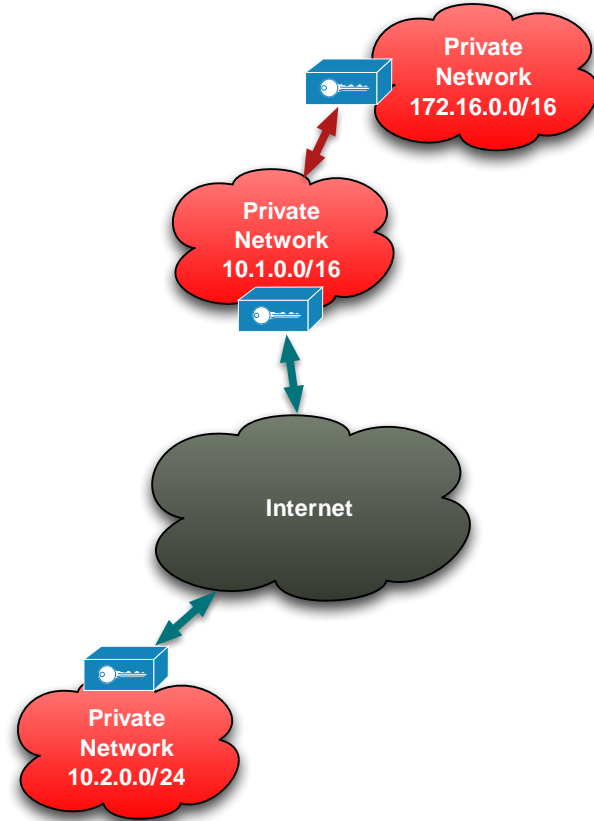
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- Public & private IP address ranges (IPv4 or IPv6)



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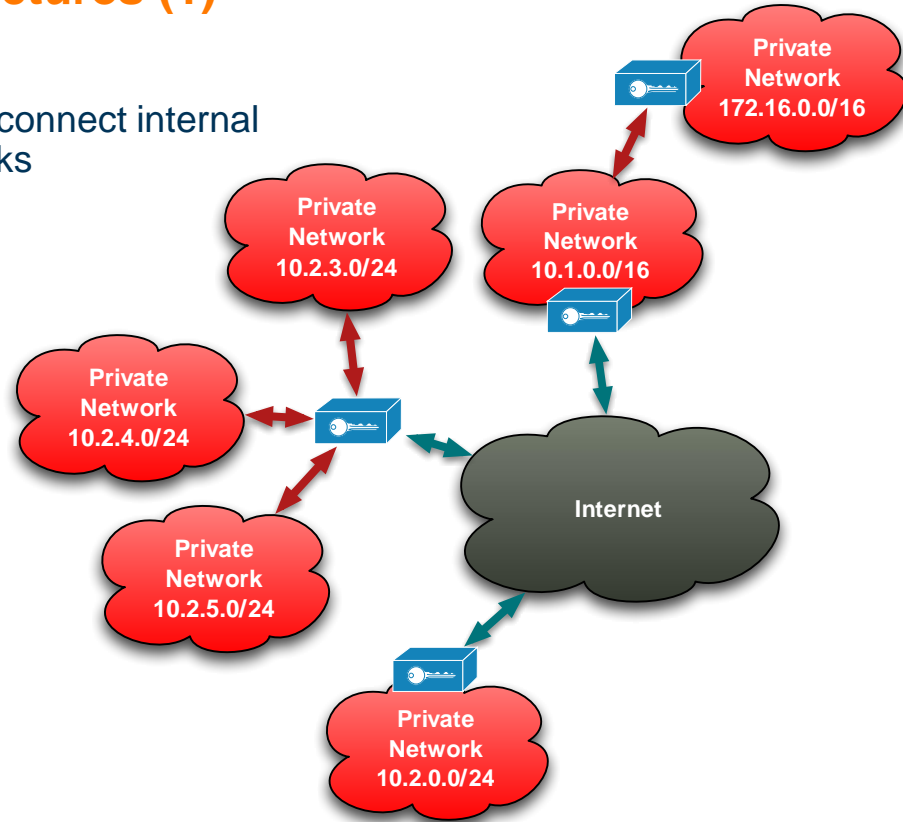
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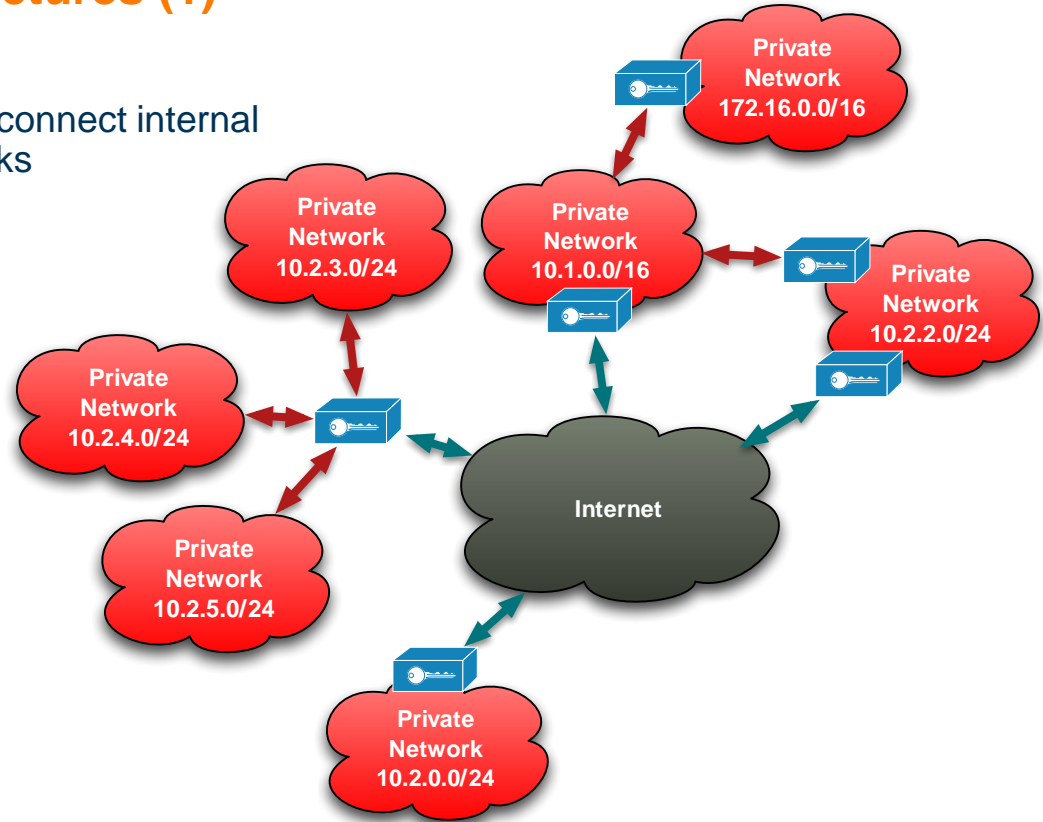
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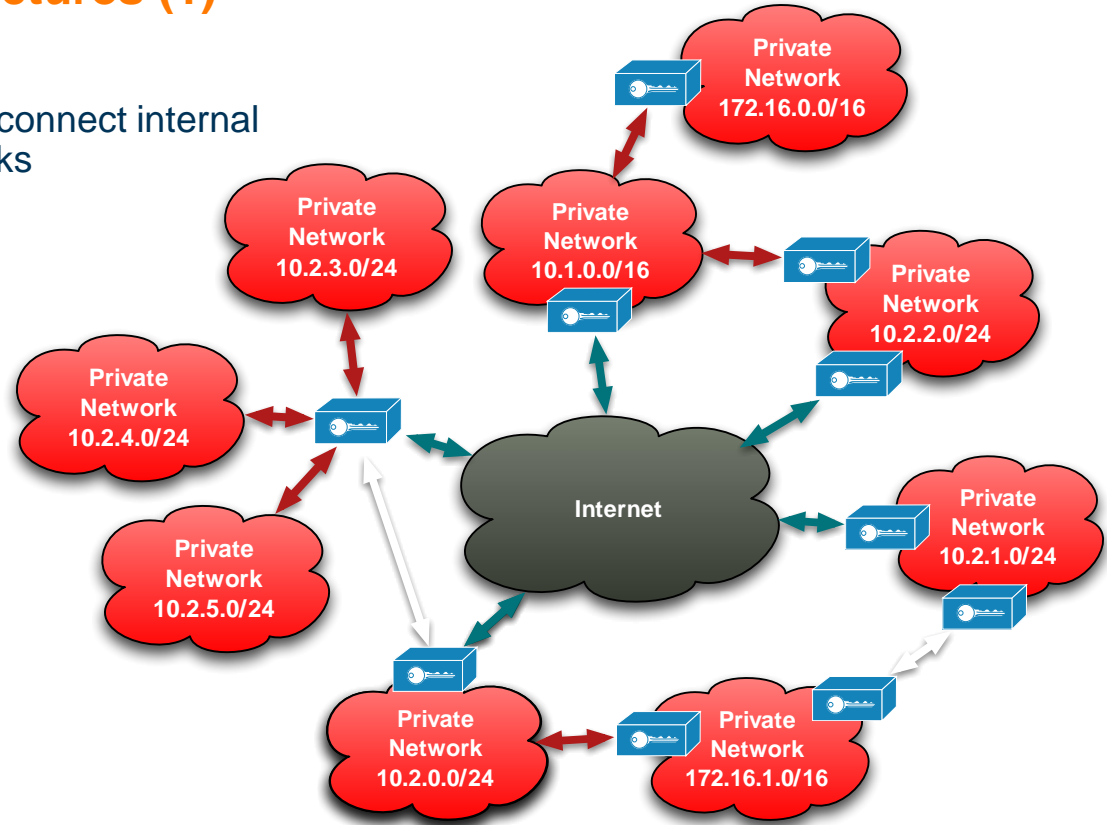
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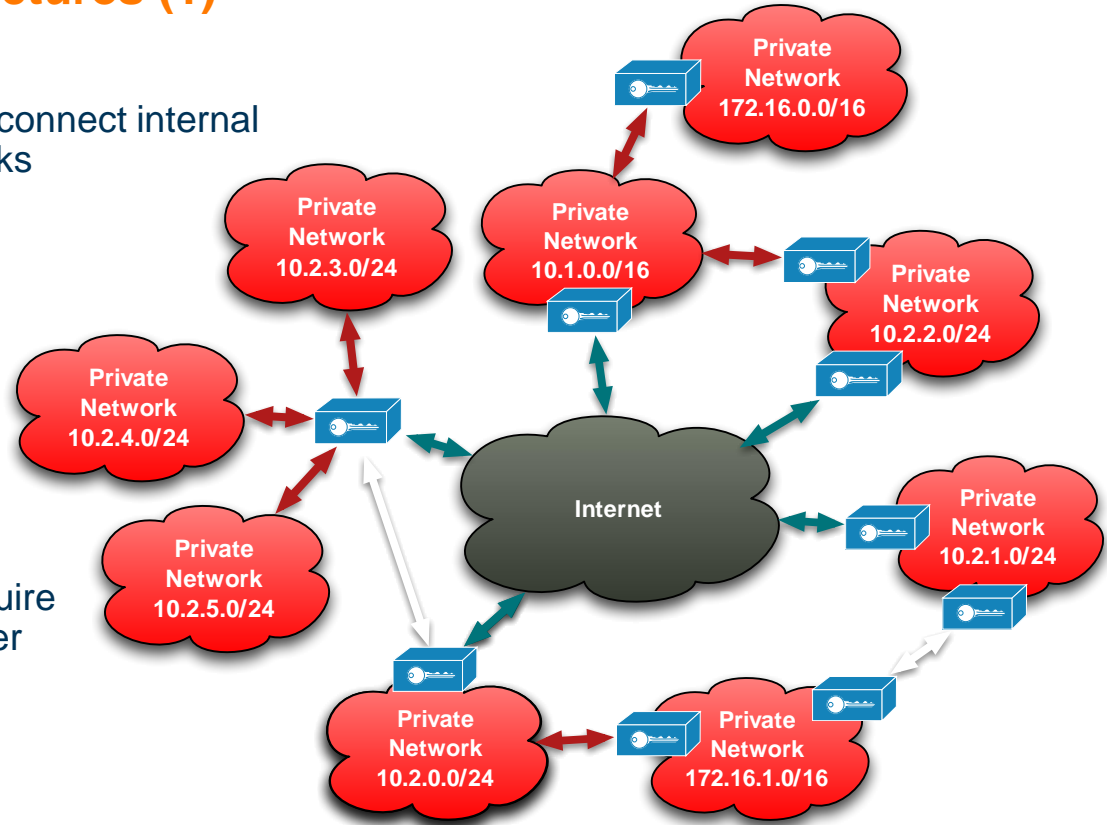
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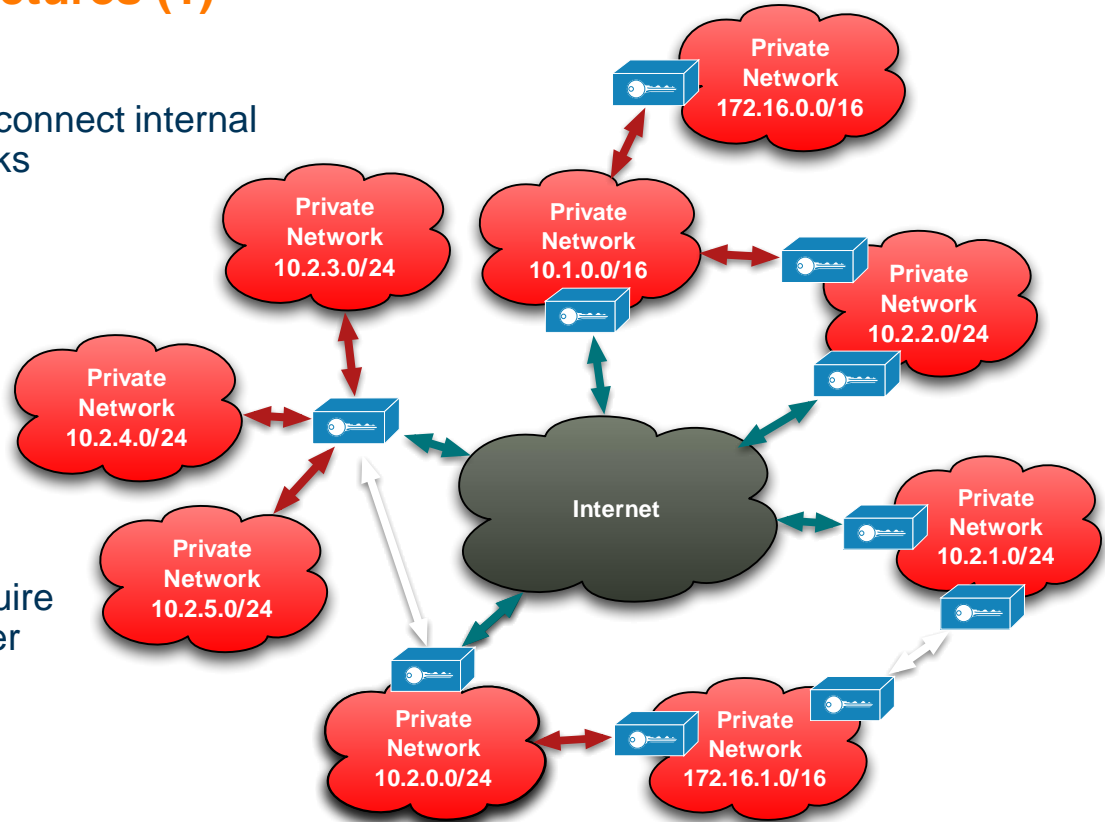


# Network Layer VPN Infrastructures (1)

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- Cycles in the network (required for robustness and handling load!)
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⇒ High complexity!



## Network Layer VPN Infrastructures (2)

### Customer expectations are simple:

- BSI-compliant crypto-processing at line speed or at least at well-defined speeds
- Handling of appliances as good/bad as other networking equipment:  
Robustness, management, enrollment
- Behave as transparently as possible
- Important VPN properties: scalability, agility, robustness

### Key enablers to implement secure, scalable and robust VPNs:

- Avoid centralized components
- Use as few security associations (SAs) as possible (SA establishment is expensive!)
  - VPN gateways implement an **overlay network/graph** (gateway = node, SA = link)
  - Use tunneled SAs to guarantee **end-to-end security** (some gateways might be compromised)
- Keep (overlay) topology knowledge local
- Automatic configuration as far as possible (by “control algorithm”)

## Network Layer VPN Infrastructures (3)

**Further required for scenarios with enhanced needs for protection (e.g., “GEHEIM”):**

- Security hardening of components, e.g., regarding:
  - Side-channel attacks
  - Minimizing trusted computing base (TCB)
  - Tamper-proofing
- Approval according to protection profile(s)

# How to Overcome “Quantum Threat” to Classical Asymmetric Cryptography?

## Two main directions:

- Post Quantum Cryptography (PQC)
  - Requires: Longer keys, longer messages and more computation (→ smart cards?)
  - Still raises concerns regarding maturity with respect to cryptanalysis
  - Required in the long run, but maybe not yet ready to be used alone
- Quantum Key Distribution (QKD)
  - Can “physically” guarantee confidentiality (but only after out-of-band authentication!)
  - Works only over “direct” medium (fiber, air) within limited reach (~100 km)

⇒ **Requires concepts for networking QKD-enabled devices**

## Open challenges:

- How to do this without unnecessarily “reinventing wheels” (→ established VPN technology)?
- How to reduce efforts for hardening “QKD networking”-related software components?
- How can security be increased between red networks with no direct QKD link?
- How can security be increased for red networks with no QKD link at all?

# Implications for QKD Integration

## At first glance: None!

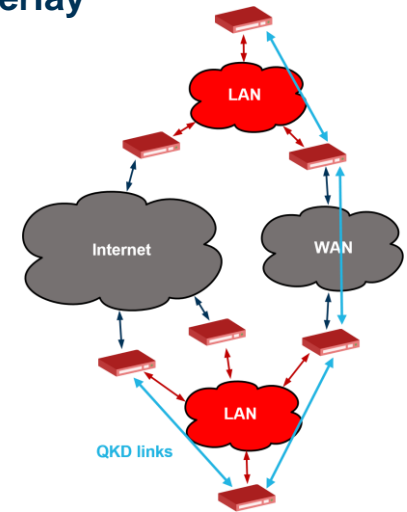
- QKD only affects confidentiality, integrity, availability properties of certain links
- “Buried” in layers below

## At second glance: We need to use the keys for establishing SAs in the overlay

- Secure interface between QKD devices and VPN gateways required
- Preferably integration of QKD keys in established protocols, e.g., IKEv2 (instead of a complete redesign)

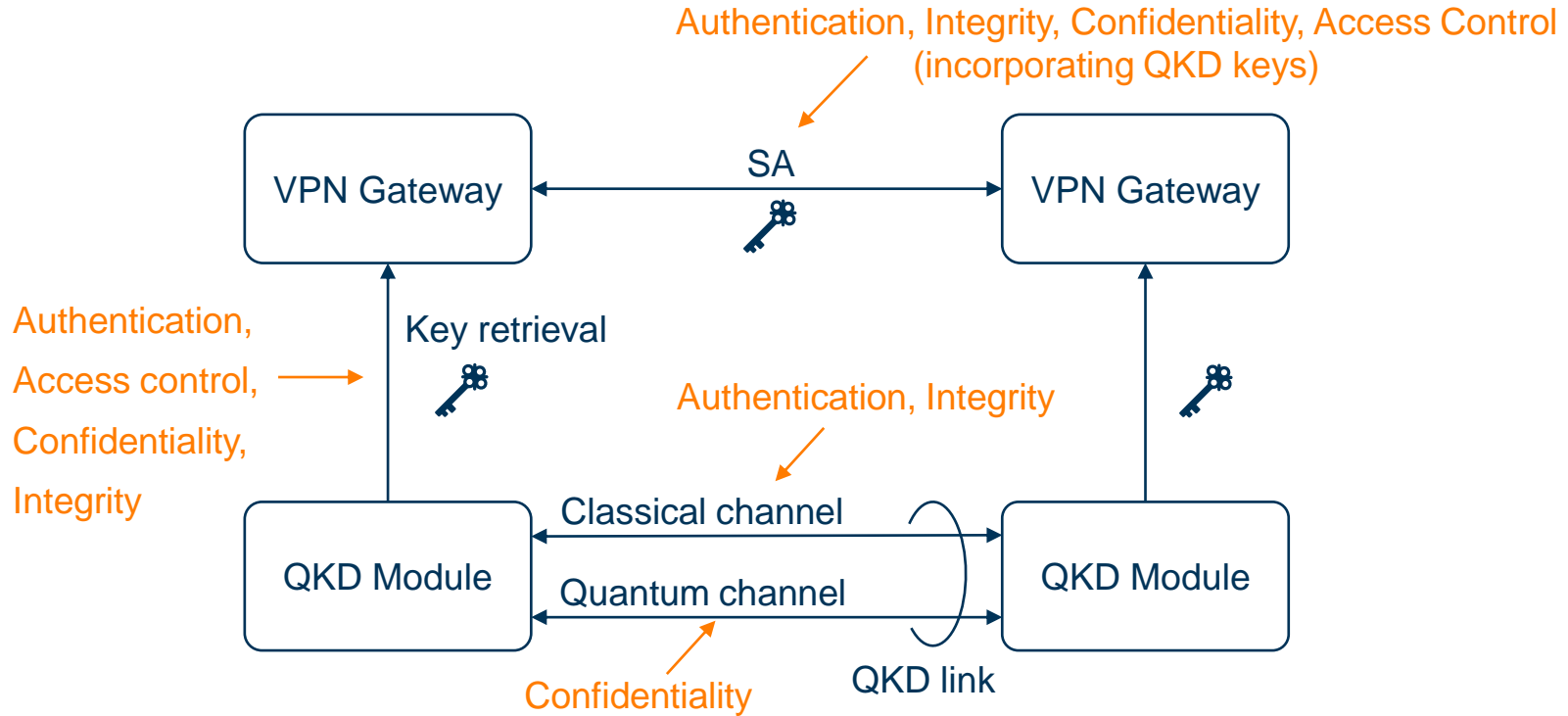
## Broader view: Impact in heterogenous infrastructures?

- Only some links have QKD (due to limited reach, costs)
- Benefit of QKD to the overall security for **arbitrary SAs**?
  - How to quantify benefit?
  - How to maximize benefit?

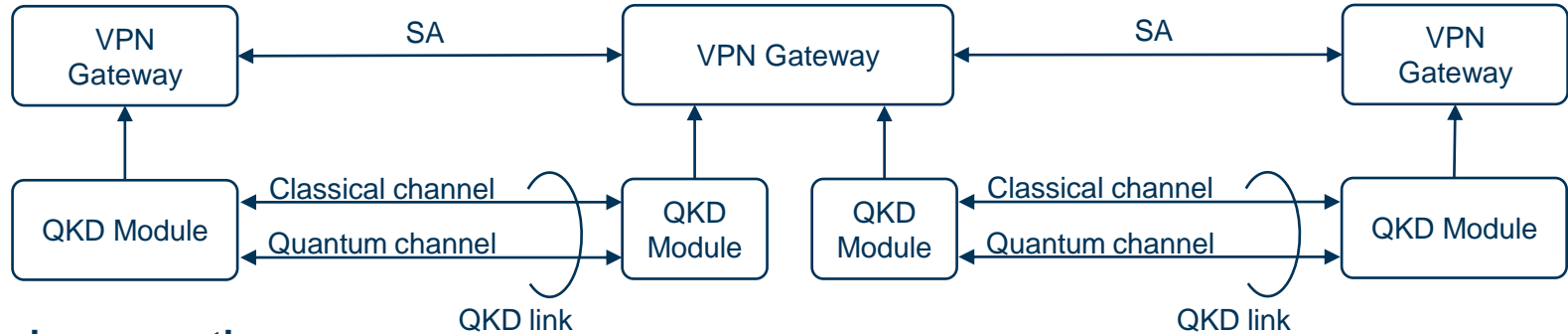


# Required Security Services

Abstract, high-level view of QKD link integration:



# QKD Network Security Considerations



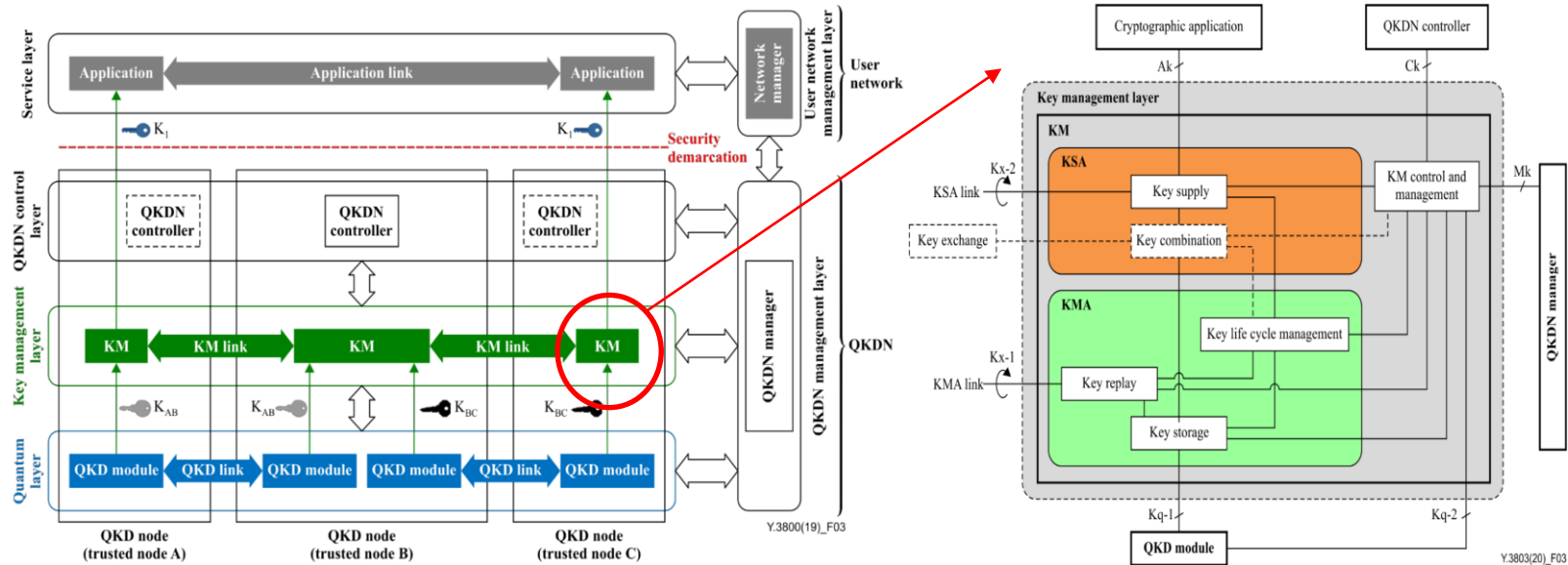
## Basic assumptions:

- Authentication needs to be realized with combination of PQC and classical cryptography
- Symmetric cryptography with sufficiently long keys (e.g.,  $\geq 256$  bit) can not be broken
- It is impossible to eavesdrop on a “securely” authenticated QKD link
- It is rather easy to eavesdrop on individual classical links
- With growing network size, it gets harder to always eavesdrop on all classical links
- It is not impossible to compromise individual VPN gateways / QKD modules (but high effort!)
- The more complex a solution is, the easier it is to compromise

# Emerging Standards: ITU-T Y.3800 – Y.3805 1/2

## Scope: QKD networks

- Idea: Transparently extend the reach of QKD by relaying keys via “trusted” nodes
- Main contribution: Reference architecture(s)





# Emerging Standards: ITU-T Y.3800 – Y.3805 2/2

## Discussion:

- No specific protocols → No interoperability, implementation complexity “hidden”
- “Standard Writer’s Standard”?
  - ~36 Functional Requirements with 9 notes [ITU-T Y.3801]
  - ~32 Functional Elements, ~22 Reference Points [ITU-T Y.3802]
  - > 50 Functions [ITU-T Y.3804]→ overly complicated?
- Security services: Identified, but very little information provided on what concrete security objectives need to be ensured and how this is supposed to be realized:
  - “[Security] [d]etails are outside the scope of this Recommendation” [ITU-T Y.3801, Y.3802, Y.3804, Y.3805]
  - “[...] security requirements described in [ITU-T X.1710], [ITU-T Y.3801] and [ITU-T Y.3802] and general network security requirements and mechanisms in IP-based networks described in [ITU-T Y.2701] and [ITU-T Y.3101] are recommended to be applied”→ How to ensure secure implementations with these recommendations?

# Gall's Law

## John Gall (1925 –2014), pediatrician and author

- Most famous book: *“General Systemantics: An Essay On How Systems Work, And Especially How They Fail...”* (1975)  
(Third edition, entitled *“The Systems Bible”* published in 2002)
- *“A complex system that works is invariably found to have evolved from a simple system that worked.  
A complex system designed from scratch never works and cannot be patched up to make it work.  
You have to start over with a working simple system.”* (1975, p. 71)
- In security, we are not only concerned with systems simply “working”, but to ensure that they do not have unintended vulnerabilities
  - This is even harder to achieve!



# Emerging Standards: ETSI GS QKD 004, 014

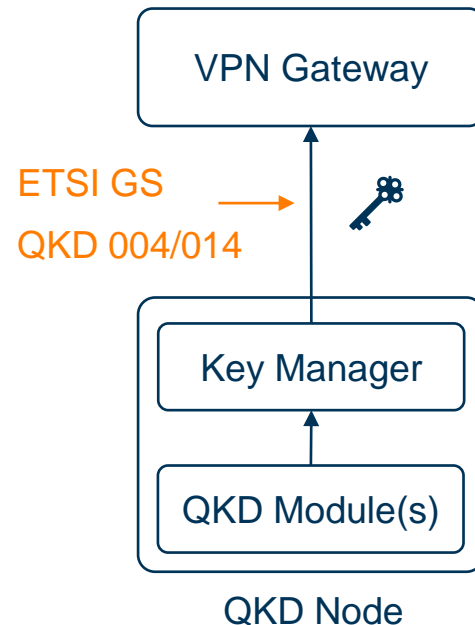
## Scope: Key retrieval in QKD networks

### ETSI GS QKD 014

- State of the art in commercially available products
- REST-based HTTP API
- Security services implemented by PKI-based TLS
  - Does not match the security level of QKD
  - Overall huge TCB: ~500k lines of code dependencies for client and server each (using well established Rust libraries)

### ETSI GS QKD 004

- Sleeker design compared to ETSI GS QKD 014 → Right direction
- But: Underspecified (e.g., encoding on wire) → Interoperability?



# Emerging Standards: ETSI GS QKD 015, 018

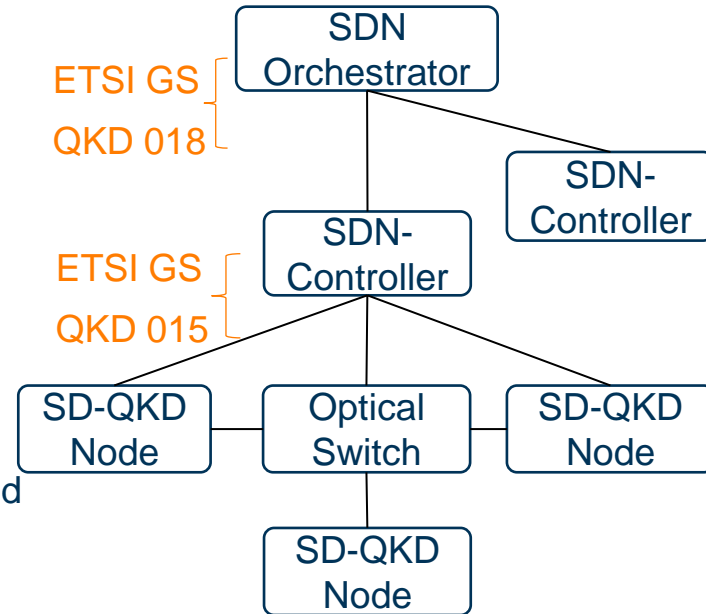
## Scope: Management & monitoring of QKD nodes

### ETSI GS QKD 015

- Central management and on demand configuration of QKD nodes and “lightpaths” using SDN
- Dynamically configuring trusted nodes to increase reachability  
→ Introduces central weak point (SDN controller)

### ETSI GS QKD 018

- Introduces SDN orchestrator for multi-domain management and monitoring
  - But: What is a domain? How separated?



## Emerging Standards: Reflection

### Common conception/objective: “Standalone” QKD networks?

- Hope: Maximizes transparency for (generic) key consumers

### Not optimally suited in the context of existing VPN infrastructures

- Routing, key management, authentication, and integrity implemented on two layers (QKD and VPN) → Increased complexity and larger TCB
- Lack of standardization for many interfaces and implementation of security services (e.g., authentication on classical channel of QKD links)
  - Proprietary protocols and implementations
  - Additional effort for hardening and approval of QKD nodes software components
- “Trusted” nodes not satisfying (or even prohibitive?) in VPNs with enhanced needs for protection

### Integrated approach better suited?

### How to maximize the benefit of QKD **without solely relying on trusted nodes?**

# Excursion: Software Vulnerabilities

## Some examples:

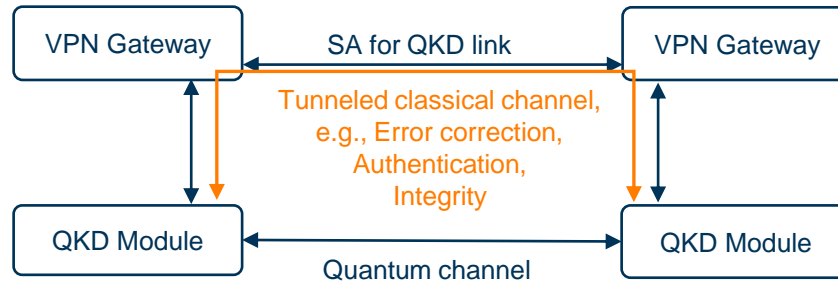
- Heartbleed [CVE-2014-01600]: Memory leak in the openssl implementation of the TLS heartbeat extension → Potentially leaked many long-term secret keys
- Log4Shell [CVE-2021-44228]: Vulnerability in “harmless” dependency (logging framework) → Allowed remote code execution for nearly ten years
- And countless more

## Implications:

- Avoid (designing and) implementing complex protocols from scratch
- Keep TCB as small as possible

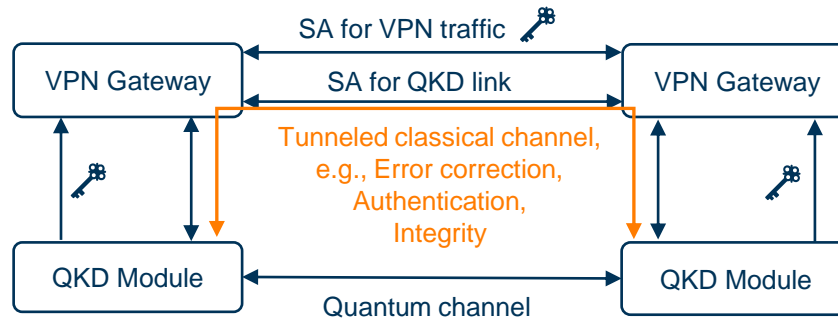
## Integrated Approach: Direct QKD Link

- Additional SA for each QKD link, established using PQC/pre-shared keys (PSKs)
- Tunnel classical channel (e.g., error correction) via VPN gateways and additional SA
- Options for security services between QKD module and VPN gateway: PQC, PSKs, “physical means”



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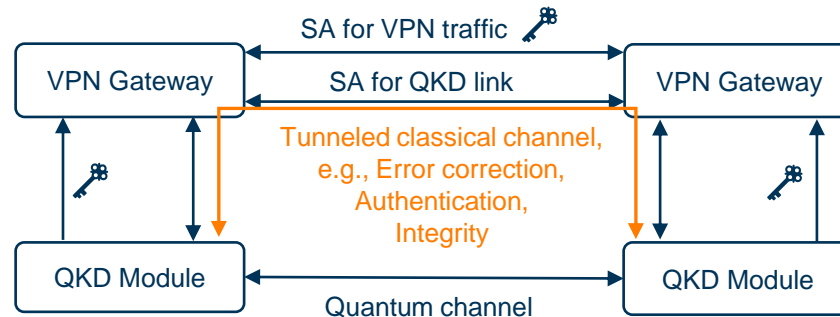
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- Traffic secured by symmetric cryptography as usual (e.g., AES, ...)





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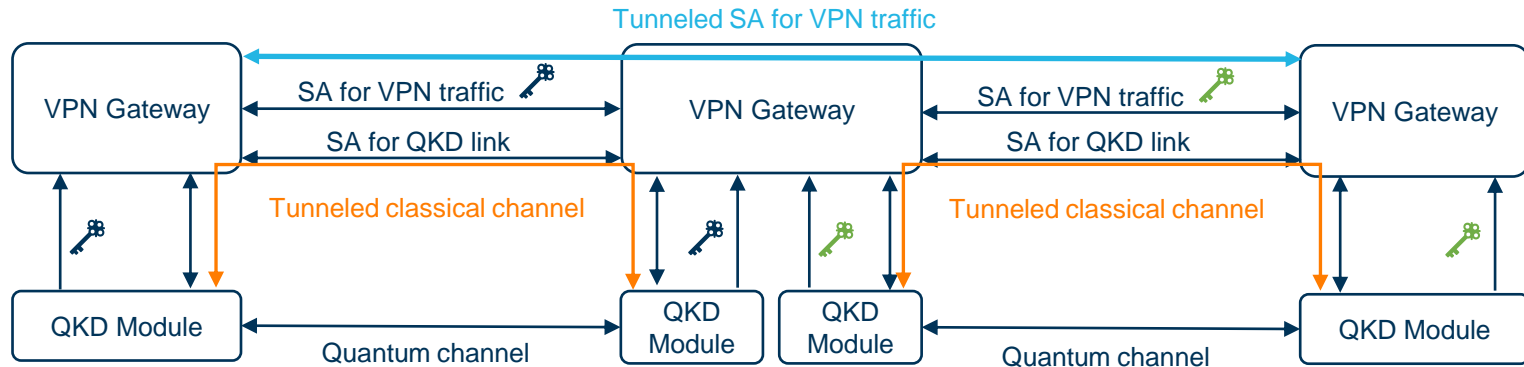


- Reduced attack surface on QKD modules (no classical communication via public channels)
- Reduced complexity of QKD modules (no authentication with other modules)

# Integrated Approach: Multi-hop QKD

## Approach:

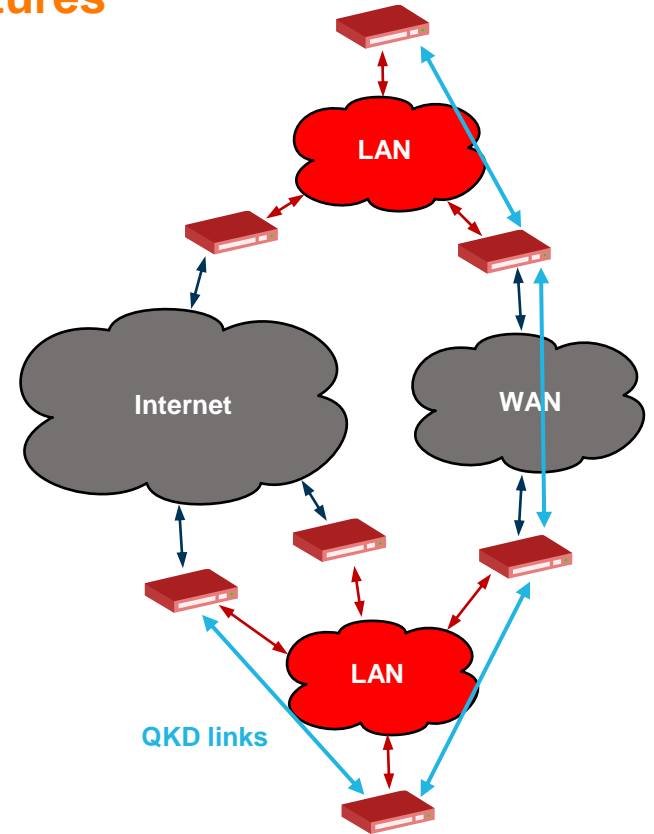
- Establish “tunneled” SA, hop-by-hop protected by existing SAs with direct access to QKD links
- End-to-end authentication and key exchange: PQC/classical cryptography
- Optimization: Re-route (shortcut) VPN traffic after successful authenticated key exchange



## Discussion:

- Same (or better?) end-to-end security properties compared to QKD network with trusted nodes
- Reduced complexity and TCB (use established VPN technologies for multi-hop key management)

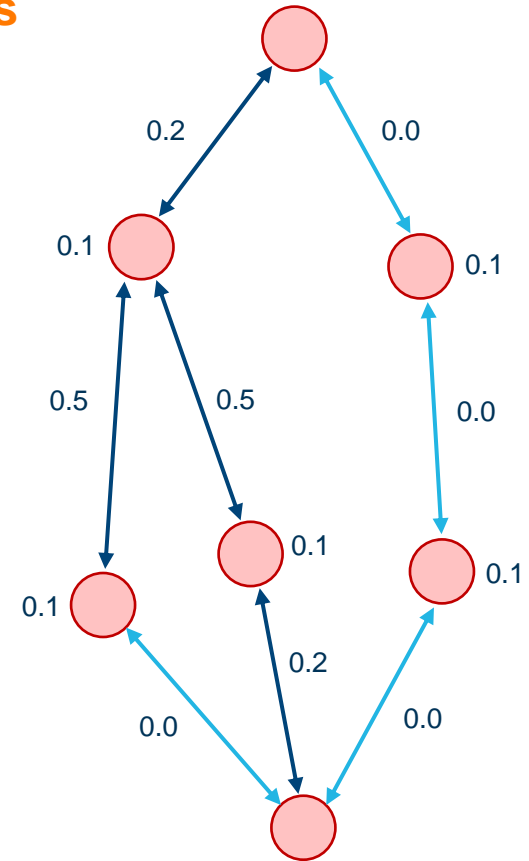
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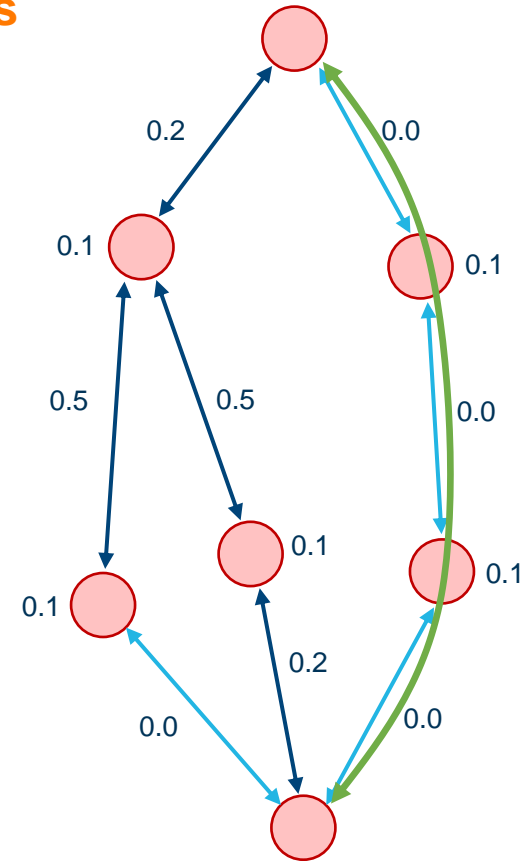
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- Augment probability of node/edge compromise



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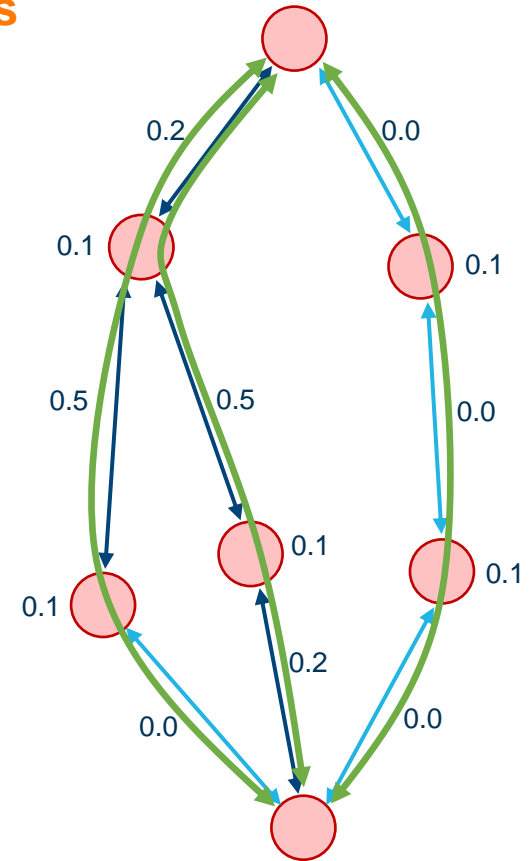
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- Graph representation of existing VPN overlay topology
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- Establish tunneled SA, e.g., on path with lowest probability of compromise
- Reinforce key by additional key exchanges
  - E.g.,  $K = H(S_1, S_2, \dots, S_n)$
  - Key shares  $S_i$  established over diverse paths and at various times



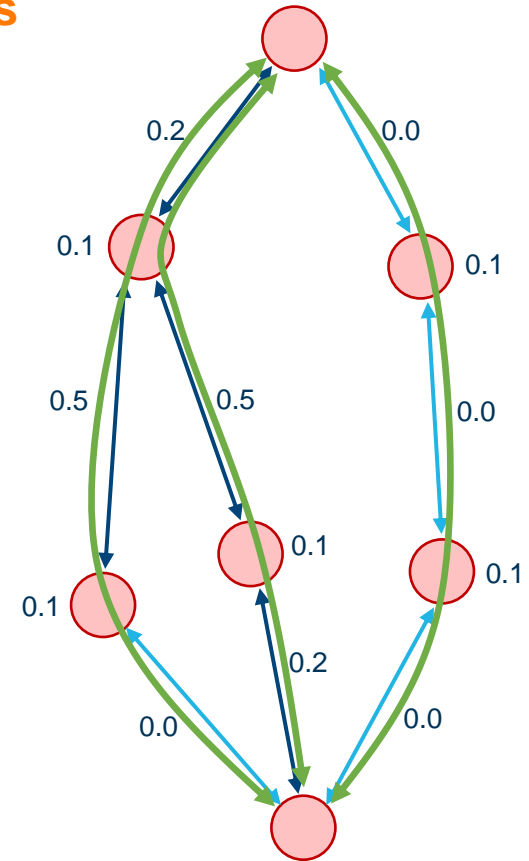
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## Discussion:

- Increases effort for attackers (must attack at all paths/times)
- Reduces dependence on “trusted” nodes
- Model is first step towards quantifying QKD gain in heterogeneous infrastructures

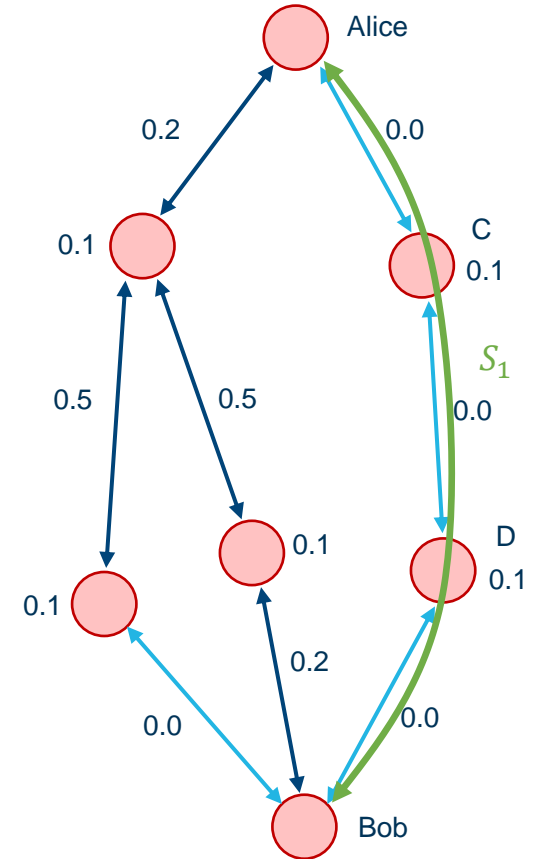


# Quantifying the Security of Keys (1)

## Scenario 1: Multi-hop QKD

- Establish key  $K_1 := S_1$  via multi-hop QKD path

$$\begin{aligned} Pr(K_1 \text{ secure}) &= Pr(C \text{ secure}) \times Pr(D \text{ secure}) \\ &= (1 - Pr(C \text{ insecure})) \times (1 - Pr(D \text{ insecure})) \\ &= (1 - 0.1) \times (1 - 0.1) = 0.81 \end{aligned}$$





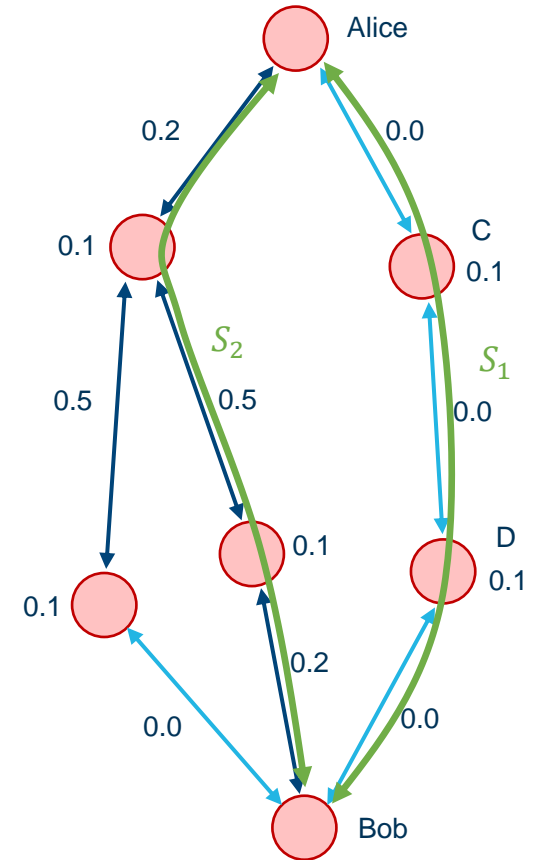
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- Reinforce key with  $S_2 \rightarrow K_2 := H(K_1, S_2)$
- $Pr(K_2 \text{ secure}) = 1 - Pr(S_1 \text{ insecure}) \times Pr(S_2 \text{ insecure})$   
 $= 1 - (1 - 0.81) \times (1 - 0.8 \times 0.9 \times 0.5 \times 0.9 \times 0.8)$   
 $\approx 0.86$



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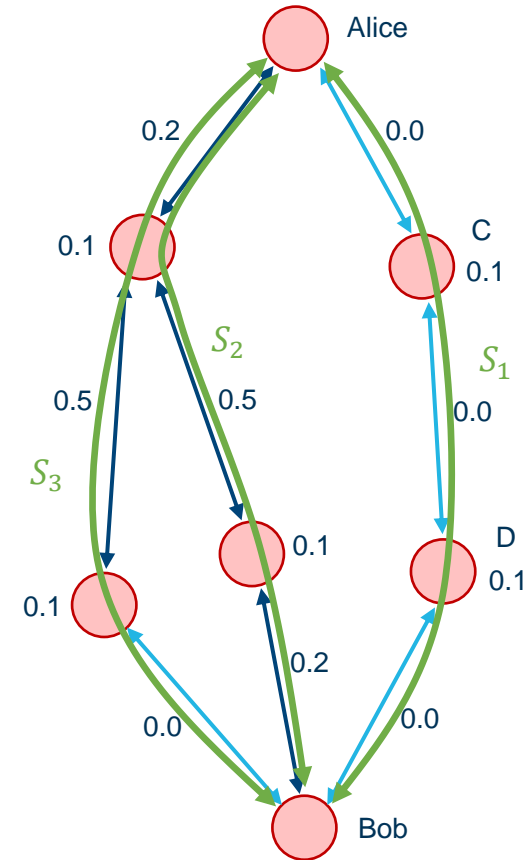
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- Further reinforce key with  $S_3 \rightarrow K_3 := H(K_2, S_3)$

$$Pr(K_3 \text{ secure}) = \dots \approx 0.90$$

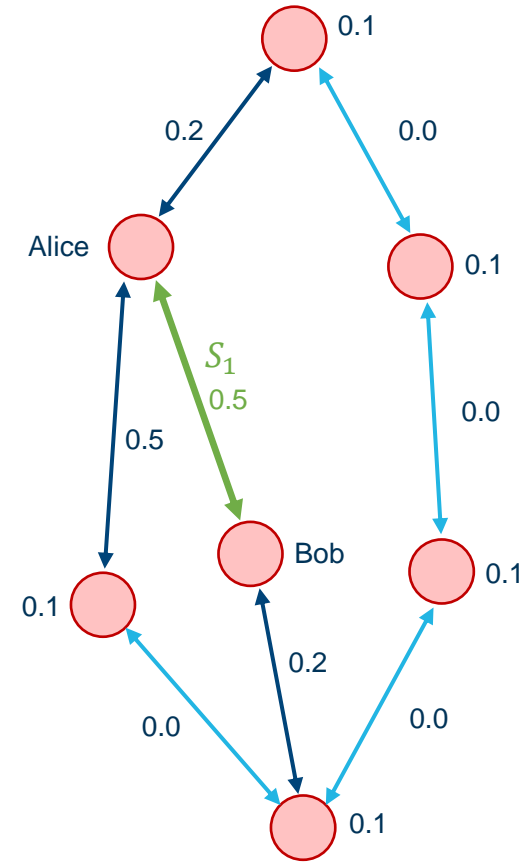
(Note: Calculation more complex due to non-disjoint paths)



## Quantifying the Security of Keys (2)

### Scenario 2: No direct access to QKD

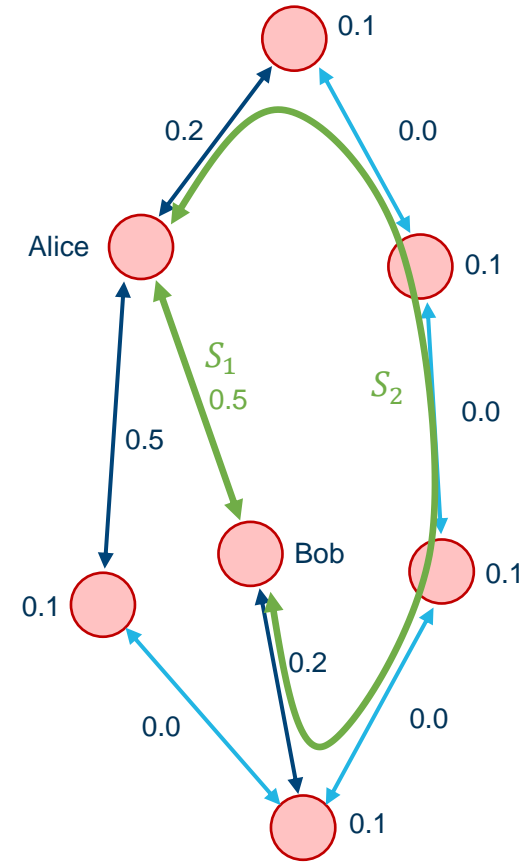
- Initial direct SA with key  $K_1 := S_1$
- $Pr(K_1 \text{ secure}) = 0.5$



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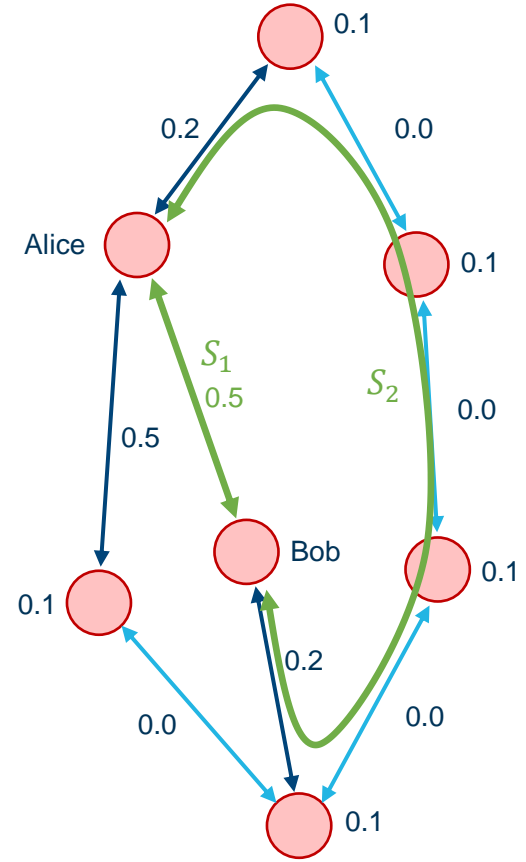
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- $S_2$  better initial choice?
- 



## Quantifying the Security of Keys (2)

### Scenario 2: No direct access to QKD

- Initial direct SA with key  $K_1 := S_1$
- $Pr(K_1 \text{ secure}) = 0.5$
  
- $S_2$  better initial choice?
- $Pr(S_2 \text{ secure}) = 0.8 \times 0.9^4 \times 0.8 \approx 0.42$   
→ No (path too long)

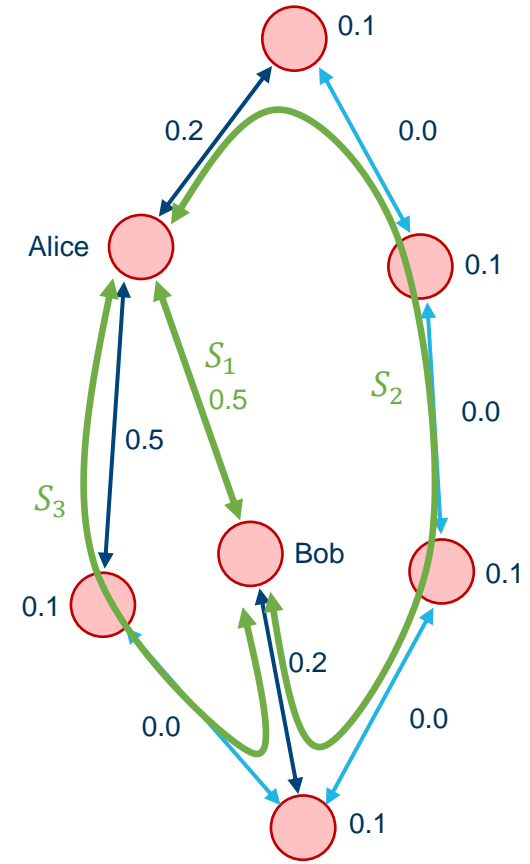




## Quantifying the Security of Keys (2)

### Scenario 2: No direct access to QKD

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- $S_2$  better initial choice?
- $Pr(S_2 \text{ secure}) = 0.8 \times 0.9^4 \times 0.8 \approx 0.42$   
→ No (path too long)
- Still: Reinforce key with  $S_2 \rightarrow K_2 := H(K_1, S_2)$
- $Pr(K_2 \text{ secure}) \approx 1 - (1 - 0.5) \times (1 - 0.42) \approx 0.71$
  
- Further reinforce key with  $S_3 \rightarrow K_3 := H(K_2, S_3)$
- $Pr(K \text{ secure}) = \dots \approx 0.79$



## Conclusion

- Security of existing VPN infrastructures threatened by quantum attackers
- One possible countermeasure: QKD
- QKD standards & commercially available products focus on “standalone” QKD networks
  - Not suited for deployment in VPN infrastructures
  - Significantly increase attack surface and TCB → Approval cumbersome
- Complexity of QKD deployment and management should be reduced by utilizing existing VPN technologies
  - Reuse existing and established VPN-Gateways to secure the classical channel between QKD devices
  - Established mechanisms for utilizing multi-hop QKD (routing, tunneled SAs)
- QKD on its own not sufficient (cost, reach) → PQC required in long-run for e2e security
- Multipath key exchange as an additional and orthogonal approach for quantum-safety



# Thanks for listening!

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